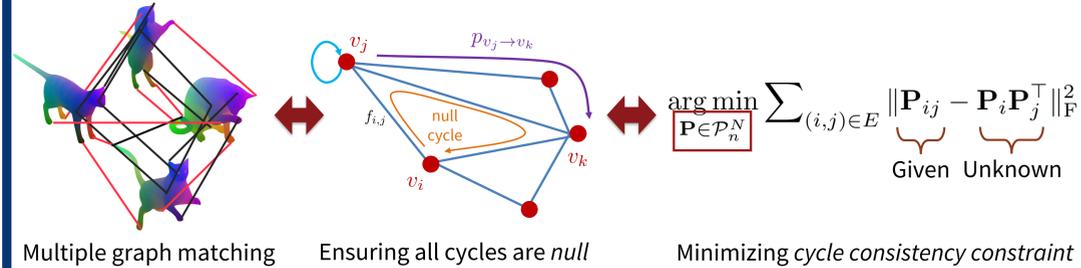


## Introduction

### Our Problem: Permutation Synchronization

Matching not just two, but  $n$  different sets of objects to each other, **jointly** [1]. In other words, a *multi-way matching*. In the scenario where correspondences are bijective, the problem converts to ensuring *cycle consistency* in the graph of **permutations** [2]:



We solve the **non-convex, combinatorial** permutation synchronization problem **without relaxation** on a **real quantum computer**.

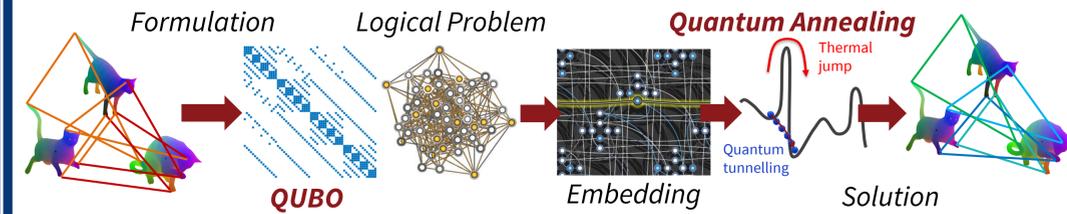
### Adiabatic Quantum Computer Vision (AQC-V)

	QUBO Suppression ECCV'20	Quantum Alignment CVPR'20	Quantum Graph Matching 3DV'2020	Quantum Sync This Paper
Quantum Processor	D-Wave 2X 1000 qubits	D-Wave 2000Q 2048 qubits	D-Wave 2000Q 2048 qubits	D-Wave Advantage 1.1 5436 qubits
AQC Problem	Non- maximum suppression	Pairwise point set alignment	Pairwise graph matching	Permutation synchronization
QPU Time	N/A	1 minute	2-3 minutes	>15 minutes

### Contributions

- (a) Formulating a **QUBO** for permutation synchronization with **permutation-ness** as a **linear** constraint
- (b) Extensive evaluation on a **real Quantum Computer, D-Wave Advantage 1.1**

## Our Approach: QuantumSync



### 1. Formulating the Vanilla QUBO

**Proposition 1.** Permutation synchronization under the Frobenius norm can be written in terms of a QUBO:

$$\arg \min_{\{X_i \in \mathcal{P}_n\}} \sum_{(i,j) \in \mathcal{E}} \|P_{ij} - X_i X_j^T\|_F^2 = \arg \min_{\{x_i \in \mathcal{P}_n\}} x^T Q' x, \quad \text{QUBO}$$

where  $x_i = \text{vec}(X_i)$ ,  $x = [\dots x_i^T \dots]^T$  and:

$$Q' = - \begin{bmatrix} I \otimes P_{11} & I \otimes P_{12} & \dots & I \otimes P_{1m} \\ I \otimes P_{21} & I \otimes P_{22} & \dots & I \otimes P_{2m} \\ \vdots & \vdots & \ddots & \vdots \\ I \otimes P_{m1} & I \otimes P_{m2} & \dots & I \otimes P_{mm} \end{bmatrix}. \quad \text{Q is symmetric and no other constraints are needed.}$$

### 2. Permutations as Linear Constraints

Group of permutations :  $\mathcal{P}_n := \{P \in \{0, 1\}^{n \times n} : P \mathbf{1}_n = \mathbf{1}_n, \mathbf{1}_n^T P = \mathbf{1}_n^T\}$

Binary                      Rows sum to 1                      Cols sum to 1

Hence, for all variables we would like to solve for :  $\text{diag}(A_1, \dots, A_n)x = \mathbf{1}, \quad A_i = \begin{bmatrix} I \otimes \mathbf{1}^T \\ \mathbf{1}^T \otimes I \end{bmatrix}$

### 3. Incorporating Linear Constraints into QUBO

**Proposition 2.** The constrained problem can be formulated into an unconstrained QUBO:

$$\arg \min_{x \in \mathcal{B}^n} x^T Q x + s^T x \quad \text{QUBO}$$

Binary Variables                      Quadratic Term                      Linear Term

$$Q = Q' + \lambda A^T A \quad s = -2\lambda A^T b$$

## Experimental Evaluation

### 1. Solving Real Problems on D-Wave Advantage 1.1

**D-Wave Python API:** [docs.ocean.dwavesys.com/en/stable/](https://docs.ocean.dwavesys.com/en/stable/)

Willow Dataset [3]

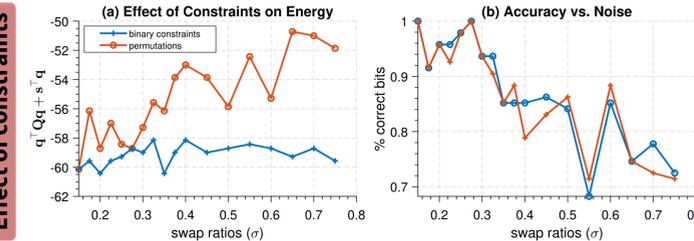


We can match the state-of-the-art methods in small problems ( $n=4, m=4$ ).

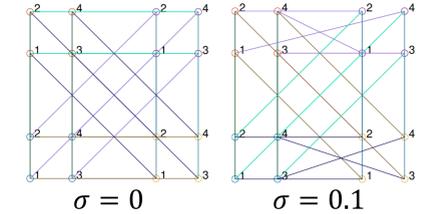
	Average
Exhaustive	<b>0.88 ± 0.104</b>
EIG	0.83 ± 0.088
ALS	0.87 ± 0.092
LIFT	0.87 ± 0.094
Birkhoff	0.87 ± 0.093
D-Wave(Ours)	0.87 ± 0.096

### 2. Impact of regularization (Binary variables vs. Permutations)

Effect of constraints

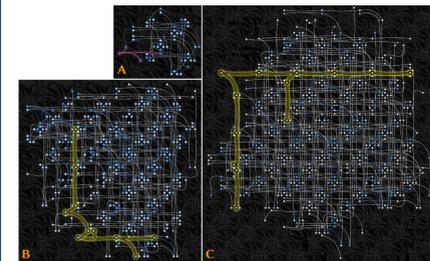


used synthetic data



Our constraint injection scheme yields valid permutations while maintaining the solution quality.

### 3. Insights into hardware implementation



We are excited for the future of quantum computers and the extent that qubits can be scaled.

For different number of points  $n$  and number of views  $m$ :  
(a) number of qubits required,  
(b) maximum chain length required on Advantage 1.1,  
(c) average number of measured optimal solutions.

- Our forward-looking experiments demonstrate that quantum hardware has reached the level that it can be applied to real-world problems.
- We hope to inspire and foster new and exciting **research in quantum computer vision**.

### References

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- Birdal, Tolga, and Umüt Simsekli. "Probabilistic permutation synchronization using the riemannian structure of the birkhoff polytope." *CVPR* 2019.
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